

Running Example: Dumbbell Code

$$H := \begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

$$C = \{ (0, 0, 0, 0, 0, 0, 0), (1, 1, 1, 0, 0, 0, 0), \\ (0, 0, 0, 0, 1, 1, 1), (1, 1, 1, 0, 1, 1, 1) \}$$

Parity Polytope of Dumbbell Code

$$H := \begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

$\Omega(H)$ is all points of the form $(a, a, a, b, c, c, c) \in \mathbb{R}^7$ with

$$0 \leq a \leq 1 \quad 0 \leq b \leq 1 \quad 0 \leq c \leq 1$$

and

$$b \leq 2a$$

$$b \leq 2c$$

$$2a + b \leq 2$$

$$b + 2c \leq 2$$

i.e.,

$$\frac{b}{2} \leq a \leq 1 - \frac{b}{2}$$

$$\frac{b}{2} \leq c \leq 1 - \frac{b}{2}$$

Vertices of $\Omega(H)$

$$\Omega(H) = \left\{ (a, a, a, b, c, c, c) \in \mathbb{R}^7 \mid \begin{array}{l} 0 \leq a \leq 1 \\ 0 \leq b \leq 1 \\ 0 \leq c \leq 1 \\ \frac{b}{2} \leq a \leq 1 - \frac{b}{2} \\ \frac{b}{2} \leq c \leq 1 - \frac{b}{2} \end{array} \right\}$$

There are five vertices:

$$(0, 0, 0, 0, 0, 0, 0)$$

$$(1, 1, 1, 0, 0, 0, 0)$$

$$(0, 0, 0, 0, 1, 1, 1)$$

$$(1, 1, 1, 0, 1, 1, 1)$$

$$\left(\frac{1}{2}, \frac{1}{2}, \frac{1}{2}, 1, \frac{1}{2}, \frac{1}{2}, \frac{1}{2}\right)$$

Cost Function Example

Suppose $\tilde{\mathbf{y}} = (0, 1, 0, 1, 0, 1, 0)$ is received (and $\mathbf{0}$ was sent) from a BSC with probability of error 0.1. Then

$$\begin{aligned}\gamma &= \left(\ln \left(\frac{.9}{.1} \right), \ln \left(\frac{.1}{.9} \right), \ln \left(\frac{.9}{.1} \right), \ln \left(\frac{.1}{.9} \right), \ln \left(\frac{.9}{.1} \right), \ln \left(\frac{.1}{.9} \right), \ln \left(\frac{.9}{.1} \right) \right) \\ &= (2.198, -2.198, 2.198, -2.198, 2.198, -2.198, 2.198)\end{aligned}$$

We have $d(\tilde{\mathbf{y}}, \mathbf{0}) = 3$ and $d(\tilde{\mathbf{y}}, \mathbf{x}) > 3$ for all $\mathbf{x} \in C \setminus \{\mathbf{0}\}$, but the cost of $(\frac{1}{2}, \frac{1}{2}, \frac{1}{2}, 1, \frac{1}{2}, \frac{1}{2}, \frac{1}{2})$ is negative and we have a decoding error.

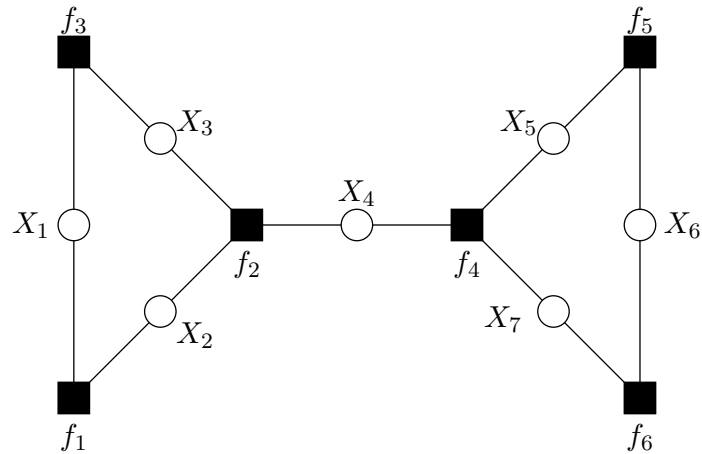
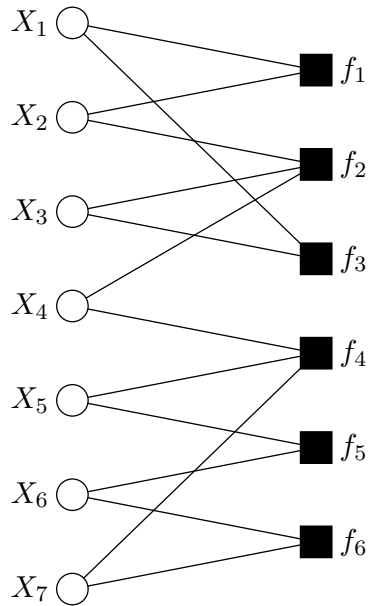
Fundamental Cone

$$H := \begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

$$\mathcal{K}(H) = \left\{ (a, a, a, b, c, c, c) \in \mathbb{R}^7 \mid \begin{array}{l} a \geq 0, b \geq 0, c \geq 0 \\ 2a \geq b, 2c \geq b \end{array} \right\}$$

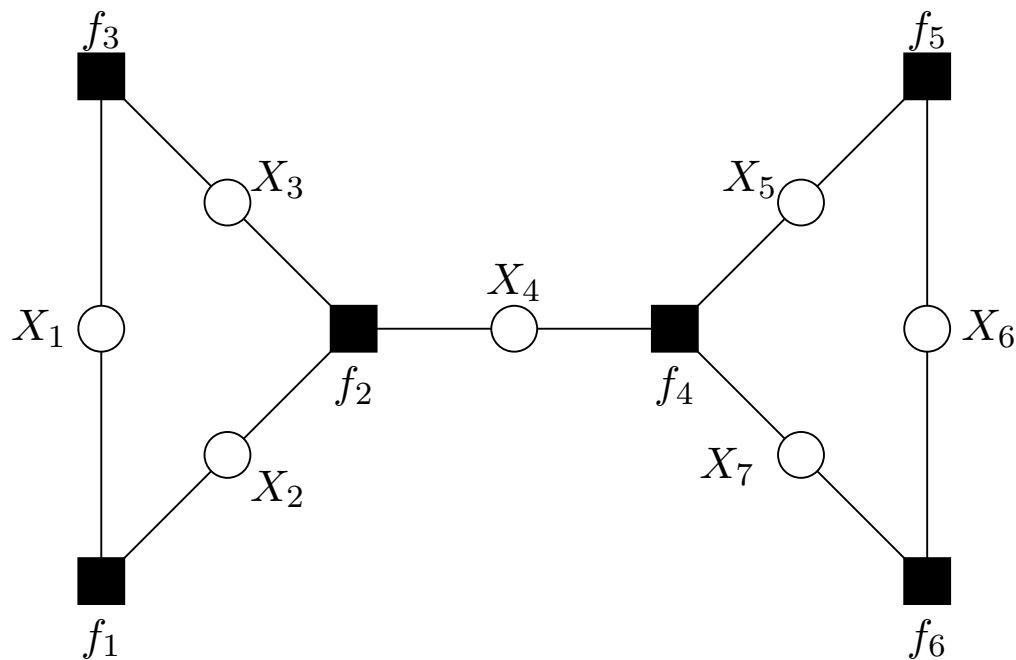
Tanner graph

$$H := \begin{pmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 1 \end{pmatrix}$$

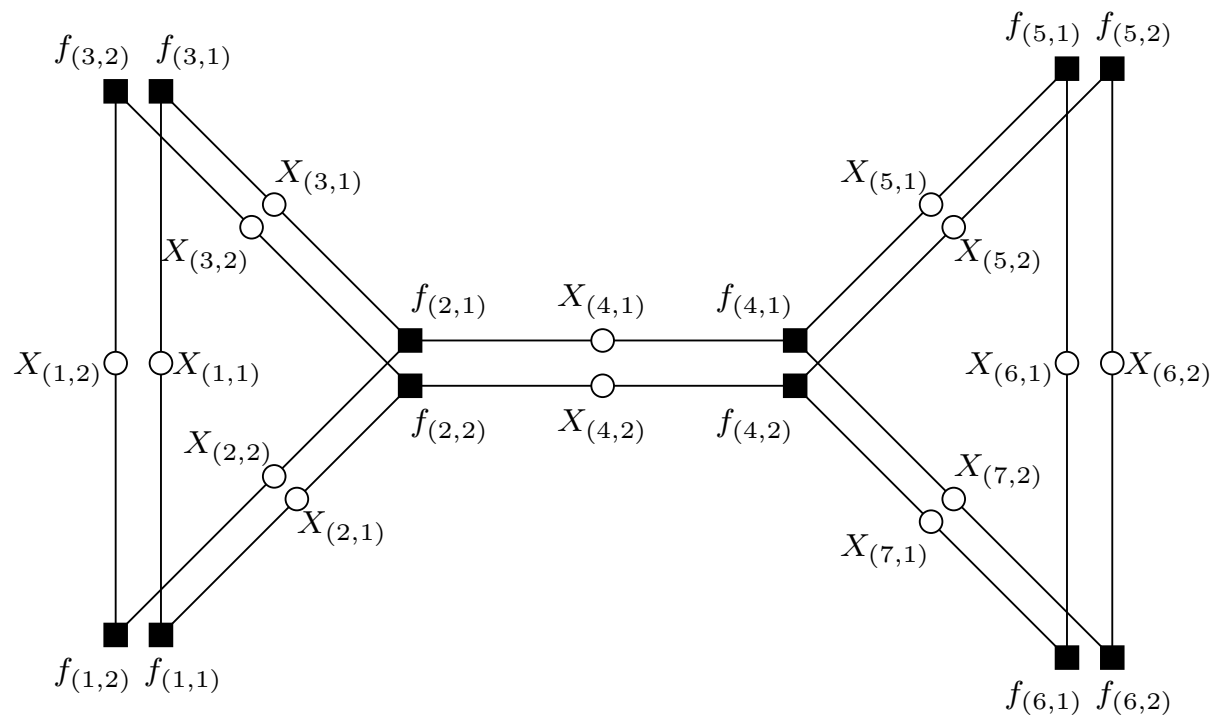


Graphical representation

$$\mathbf{c} = (1, 1, 1, 0, 0, 0, 0)$$

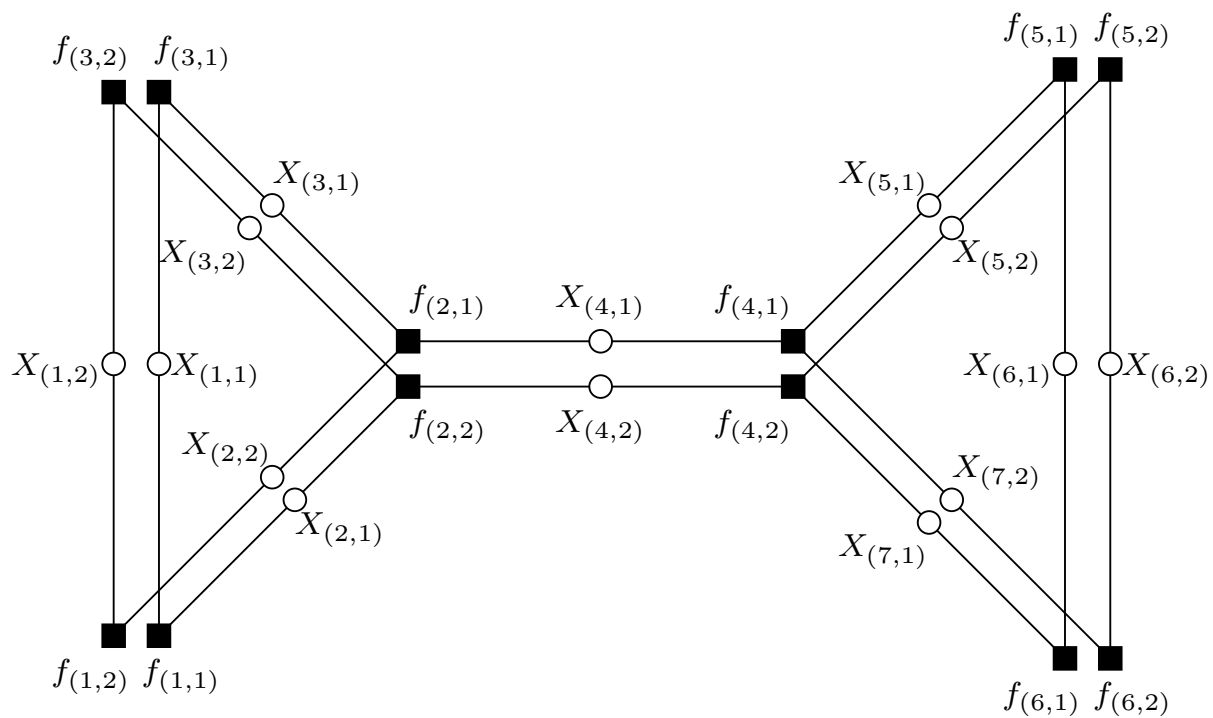


A 2-cover of the Dumbbell Graph



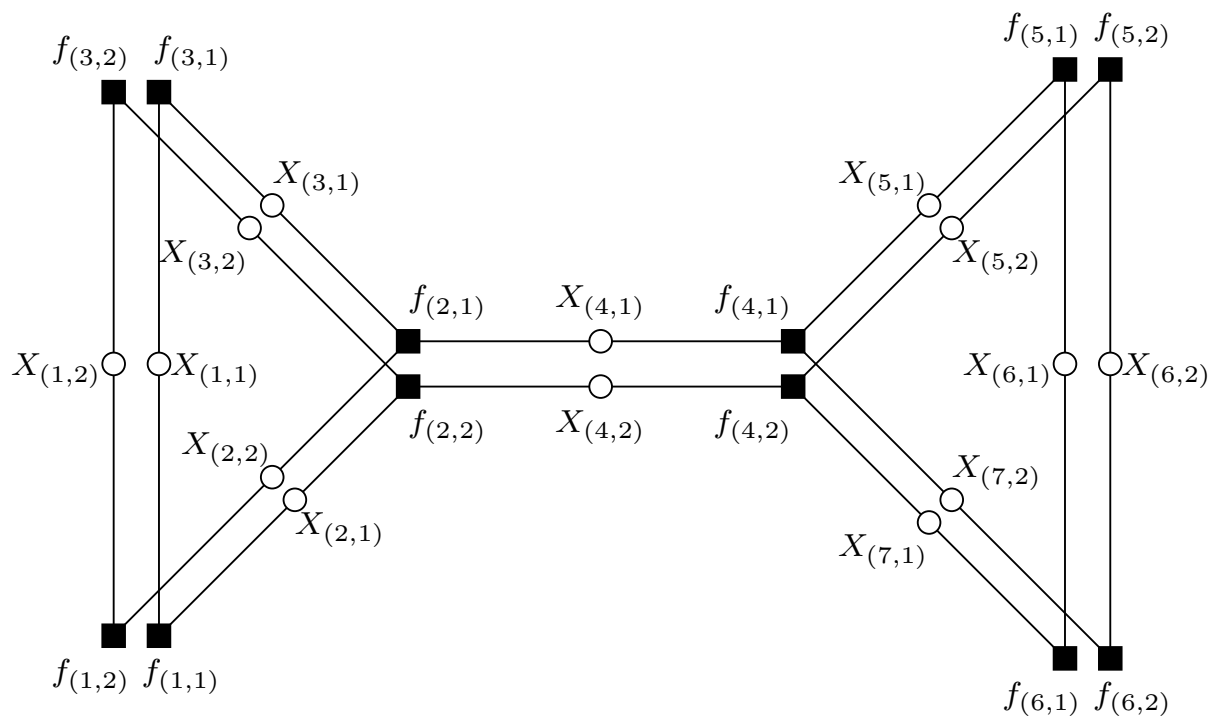
A lift to the 2-cover

$$\tilde{c} = (1:1, 1:1, 1:1, 0:0, 0:0, 0:0, 0:0)$$



A codeword on the 2-cover which is not a lift

$$\tilde{\mathbf{a}} = (1:0, 1:0, 1:0, 1:1, 1:0, 1:0, 1:0)$$



Pseudo-codewords

The pseudo-codeword associated to

$$\tilde{\mathbf{c}} = (1:1, 1:1, 1:1, 0:0, 0:0, 0:0, 0:0)$$

is

$$\mathbf{p}(\tilde{\mathbf{c}}) = (2, 2, 2, 0, 0, 0, 0).$$

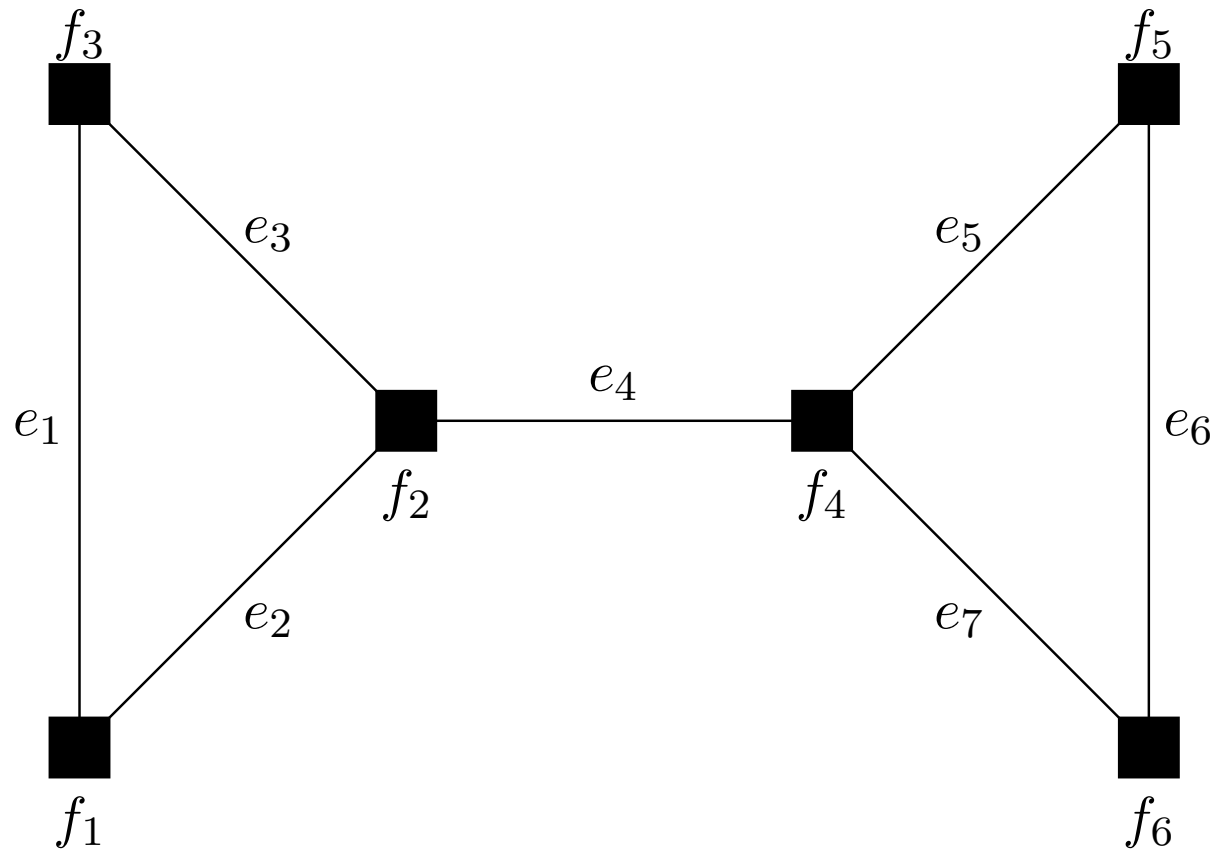
The pseudo-codeword associated to

$$\tilde{\mathbf{a}} = (1:0, 1:0, 1:0, 1:1, 1:0, 1:0, 1:0)$$

is

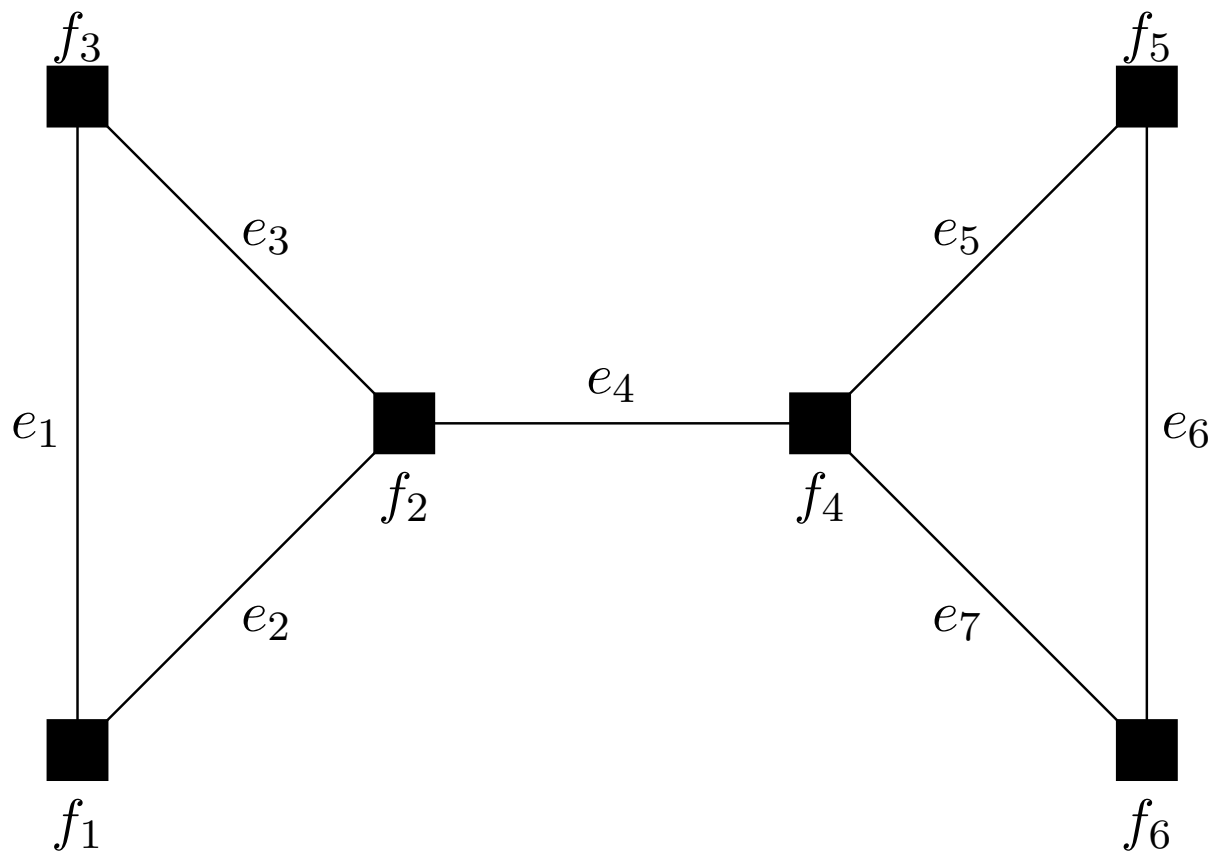
$$\mathbf{p}(\tilde{\mathbf{a}}) = (1, 1, 1, 2, 1, 1, 1).$$

Normal graph

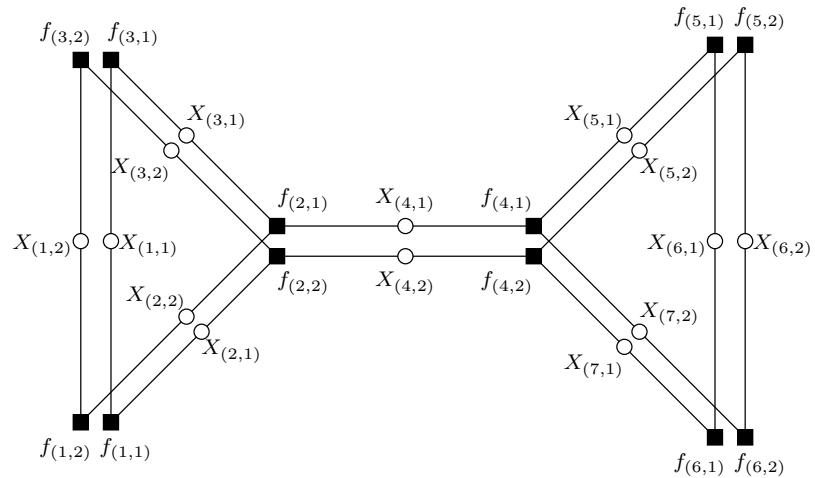
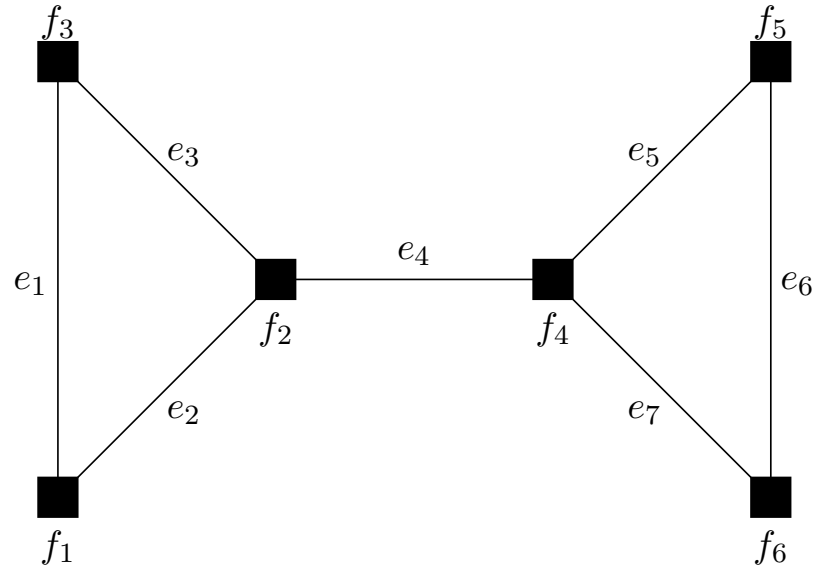


Normal graph

The codeword $(1, 1, 1, 0, 0, 0, 0)$ corresponds to the cycle (e_1, e_2, e_3) .

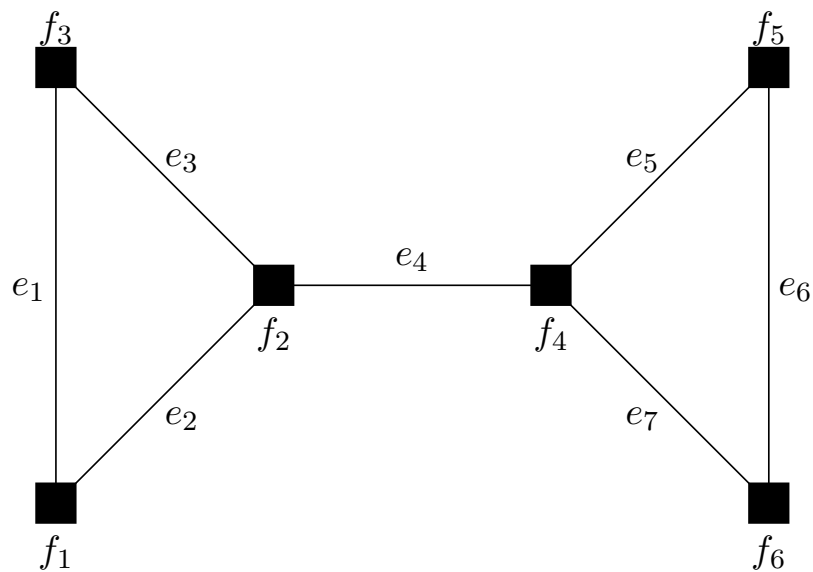


Pseudo-codeword of a cycle code



Monomial of a path

The monomial of the path $(e_1, e_2, e_4, e_5, e_6, e_7, e_4, e_3)$



is $u_1 u_2 u_3 u_4^2 u_5 u_6 u_7$.

Edge zeta function for Dumbbell Code

Let N be the normal graph for the Dumbbell Code. Then

$$\begin{aligned}\zeta_N(u_1, \dots, u_7)^{-1} &= 1 - 2u_1u_2u_3 + u_1^2u_2^2u_3^2 - 2u_5u_6u_7 \\ &\quad + 4u_1u_2u_3u_5u_6u_7 - 2u_1^2u_2^2u_3^2u_5u_6u_7 - 4u_1u_2u_3u_4^2u_5u_6u_7 \\ &\quad + 4u_1^2u_2^2u_3^2u_4^2u_5u_6u_7 + u_5^2u_6^2u_7^2 - 2u_1u_2u_3u_5^2u_6^2u_7^2 \\ &\quad + u_1^2u_2^2u_3^2u_5^2u_6^2u_7^2 + 4u_1u_2u_3u_4^2u_5^2u_6^2u_7^2 - 4u_1^2u_2^2u_3^2u_4^2u_5^2u_6^2u_7^2.\end{aligned}$$

Edge zeta function for Dumbbell Code II Expanding out the Taylor series gives

$$\begin{aligned}
\zeta_N(u_1, \dots, u_7) = & 1 + 2u_1u_2u_3 + 3u_1^2u_2^2u_3^2 + 2u_5u_6u_7 + 4u_1u_2u_3u_5u_6u_7 \\
& + 6u_1^2u_2^2u_3^2u_5u_6u_7 + 4u_1u_2u_3u_4^2u_5u_6u_7 + 12u_1^2u_2^2u_3^2u_4^2u_5u_6u_7 \\
& + 3u_5^2u_6^2u_7^2 + 6u_1u_2u_3u_5^2u_6^2u_7^2 + 9u_1^2u_2^2u_3^2u_5^2u_6^2u_7^2 \\
& + 12u_1u_2u_3u_4^2u_5^2u_6^2u_7^2 + 36u_1^2u_2^2u_3^2u_4^2u_5^2u_6^2u_7^2 + \dots .
\end{aligned}$$

More precisely ...

- Label the bit nodes of T as x_1, \dots, x_n
- Label the edges of T as $e_{(i,j)}$, $1 \leq i \leq n$ and $1 \leq j \leq d_i$, so that $e_{(i,1)}, \dots, e_{(i,d_i)}$ are incident to x_i
- Set $N := \sum d_i = \#$ of edges of T
- Define $\phi : \mathbb{F}_2^N \rightarrow \mathbb{F}_2^n$ by

$$\phi((\mathbf{c})_{(i,j)}) = (c_{(1,1)}, c_{(2,1)}, \dots, c_{(n,1)})$$

- $\widehat{C}' = \{\mathbf{c} \in \widehat{C} \mid c_{(i,1)} = \dots = c_{(i,d_i)} \text{ for all } i\}$

Then $C = \phi(\widehat{C}')$.